Persistent Data Structure

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What is?

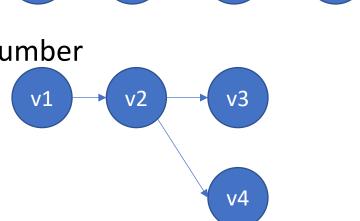
- Traditional data structure is ephemeral
 - When you modify it, it change.
 - The old data is lost, there is no obvious way to backtrack the edit (in most cases)
- Persistent Data Structure is a data structure that allow us to see back in time
 - Time is identified by version number
 - Partial persistent: can read back in time
 - Full persistent: can read/write back in time, creating fork of versions

When to use Persistent?

- Update query style problem
- Where query is NOT known beforehand
 - E.g., interactive task
- If query is known, it is better to "sweep" though time
 - E.g., sort query, update by times
- Most of the time, it is not that obvious

Example Interface

- Consider array of size n, A
- Normal array of int
 - Read: int get(idx) (e.g., cout << A[x])
 - Write: void set(idx,value) (e.g., A[x] = 20)
- Partial Persistent Array
 - Versioning: int current_version()
 - Read: int get(idx, version)
 - Write: int set(idx, value) //return version number
- Full persistent Array
 - Write: int set(idx, value, version)
 - Versioning: int previous_version(version)



Naïve Implementation

- Use 2D array to store 1D Array
- The other dimension is "version"
- Full-Copy of each version

operation	Ver	A[0]	A[1]	A[2]	A[3]	A[4]	A[5]	A[6]
Init	0	0	0	0	0	0	0	0
Set(3,1)	1	0	0	0	1	0	0	0
Set(5,4)	2	0	0	0	1	0	4	0
Set(2,3)	3	0	0	3	1	0	4	0
Set(3,5)	4	0	0	3	5	0	4	0
Set(5,9)	5	0	0	3	5	0	9	0

Naïve Implementation

- Set = O(n)
- Get = O(1)
- Space O(n) per operation
 - For m operation, it's O(nm)

• This approach can be done in most data structure

Today Topic

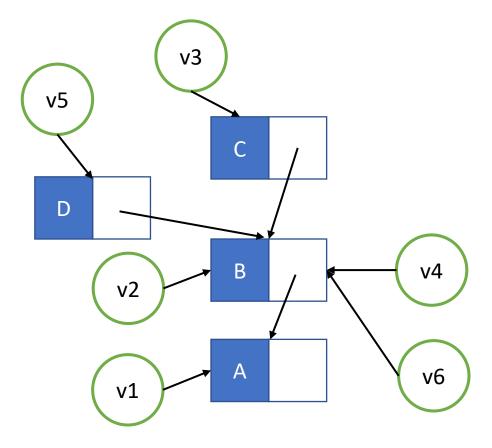
• How to make "faster" and "smaller" persistent data structure

Persistent Stack

- Idea: Pointer based (can use vector)
- Interface:
 - Create new version: push, pop
 - No change: top
- Each element in the stack contains a pair <data, next-to-top>
- Version = a pointer to the top

Example

- Sequence of operation
 - Start with version 0 = empty stack
 - Push(A) → version 1
 - Push(B) \rightarrow version 2
 - Push(C) \rightarrow version 3
 - Pop() \rightarrow version 4
 - Push(D) \rightarrow version 5
 - Pop → version 6
- O(1) push, pop, top
- O(1) space per push, pop



Code

```
Data:

vector<pair<T,int>> nodes;
vector<int> tops = {-1};
int current_version = 0;
```

```
top:

T& top(int version) {
   return nodes[tops[version]].first;
}
```

```
push:
int push(T value) {
  nodes.push_back( {value, tops.size()-1 } );
  tops.push_back(nodes.size()-1);
  return ++current_version;
}

int pop() {
  int tos = tops[current_version];
  tops.push_back(nodes[tos].second);
  return ++current_version;
}
```

Persistent Queue

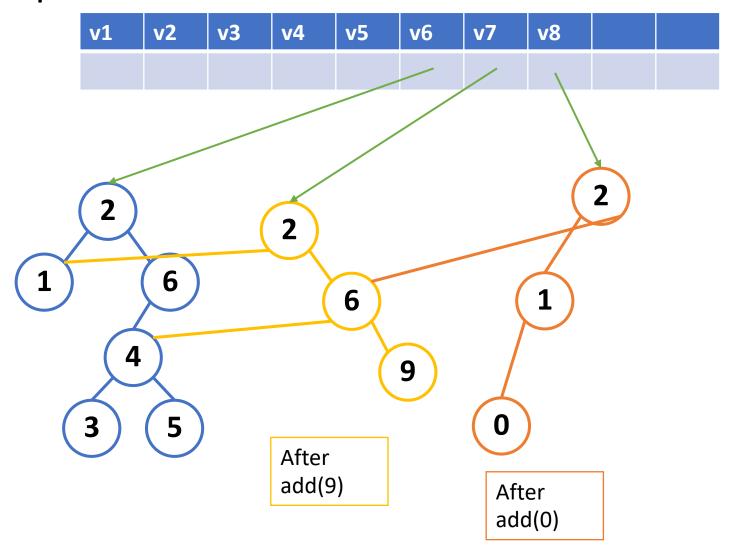
- Like a stack
- Each version keep <front,back>
- Enqueue, dequeue add new pairs of <front, back>

Persistent BS Tree

- Each version has different root
- When add, create a new version of nodes from the new leaf to the new root

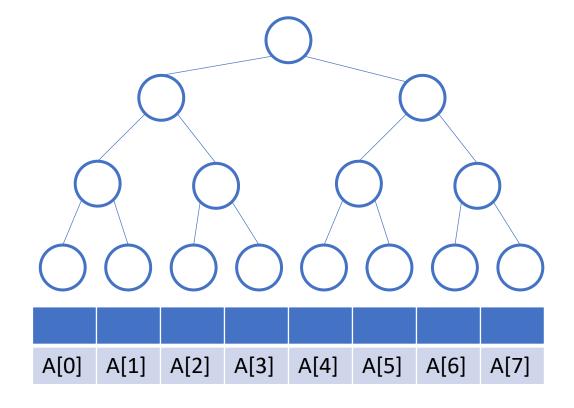
- Each insert cost additional O(h) (to the original O(h)) in times where h
 is the height of the tree
 - Also add O(h) space

Example



Persistent Array

- Use tree to store array
- Internal nodes covers some part of array (in segment tree fashion)
- Update in the same way as BST



Implementation Summary

- Queue, Stack use pointer-based
- Use "Segtree Structure" (as in Persistent Array) if underlying data structure is Array
 - Such as Array, Fenwick Tree, Priority Queue, Open Addressing Hash
 - Also Disjoint Set, Queue, Stack can be done this way
 - Need additional O(log W) where W is the size of the array
- Use "Tree Structure" (as in Persistent BST) if underlying data structure is a Tree
 - BST
 - Generalized Seg Tree

Useful Link

Link from GfG https://www.geeksforgeeks.org/persistent-data-structures/

Good Intro http://www.toves.org/books/persist/

 Purely Functional Data Structure https://www.cs.cmu.edu/~rwh/theses/okasaki.pdf

Some Interesting Problem

- CodeForces Sign on Fence http://codeforces.com/problemset/problem/484/E
- SPOJ K-th Number
- SPOJ Count on Tree
- CodeForces The Classic Problem
 - https://codeforces.com/problemset/problem/464/E